

Lecture 4

Gauss-Seidel Method (GS)

Jinn-Liang Liu
2017/4/18

We consider again Example 3.1

$$\begin{bmatrix} 3 & 2 \\ 1 & 5 \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \end{bmatrix} = \begin{bmatrix} 5 \\ 6 \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \end{bmatrix} \quad (4.1)$$

Jacobi's Method:

$$\begin{cases} x_1^{(1)} = (b_1 - a_{12}x_2^{(0)})/a_{11} \\ x_2^{(1)} = (b_2 - a_{21}x_1^{(0)})/a_{22} \end{cases} \quad (4.2)$$

For the GS method, we replace $x_1^{(0)} = 0$ by $x_1^{(1)}$, the most recent value for x_1 ,

$$\therefore x_2^{(1)} = (b_2 - a_{21}x_1^{(1)})/a_{22} \quad (4.3)$$

$$\Rightarrow \begin{cases} x_1^{(1)} = (b_1 - a_{12}x_2^{(0)})/a_{11} = (5 - 2 \cdot 0)/3 = \frac{5}{3} \\ x_2^{(2)} = (b_2 - a_{21}x_1^{(1)})/a_{22} = (6 - 1 \cdot \frac{5}{3})/5 = \frac{13}{15} \end{cases} \quad (4.4)$$

In general, if we proceed as with the Jacobi method, but now assume that the equation are examined one at a time in sequence, and that previously computed results are used as soon as they are available, we obtain the

Gauss-Seidel method

$$x_i^{(k)} = (b_i - \sum_{j < i} a_{ij}x_j^{(k)} - \sum_{j > i} a_{ij}x_j^{(k-1)})/a_{ii} \quad (4.5)$$

Table 4.1. GS Iteration						
k	0	1	2	3	\dots	∞
$x_1^{(k)}$	0	$\frac{5}{3}$	$\frac{49}{45}$	\dots	\dots	1
$x_2^{(k)}$	0	$\frac{13}{15}$	$\frac{221}{225}$	\dots	\dots	1

Considering the matrix form

$$A \vec{x} = \begin{bmatrix} a_{11} & a_{12} & \cdots & \cdots & a_{1N} \\ a_{21} & a_{22} & \cdots & \cdots & a_{2N} \\ \vdots & & \ddots & & \vdots \\ \vdots & & & \ddots & \vdots \\ a_{N1} & a_{N2} & \cdots & \cdots & a_{NN} \end{bmatrix} \begin{bmatrix} x_1 \\ x_2 \\ \vdots \\ \vdots \\ x_N \end{bmatrix} = \begin{bmatrix} b_1 \\ b_2 \\ \vdots \\ \vdots \\ b_N \end{bmatrix} \quad (4.6)$$

for the case of $i = 3$, (4.5) can be written as

$$\begin{bmatrix} \times & 0 & \cdots & \cdots & 0 \\ 0 & \times & 0 & \cdots & \vdots \\ \vdots & 0 & a_{33} & 0 & \vdots \\ \vdots & \vdots & 0 & \times & 0 \\ 0 & \cdots & \cdots & 0 & \times \end{bmatrix} \begin{bmatrix} x_1^{(k)} \\ \vdots \\ x_3^{(k)} \\ x_4^{(k-1)} \\ \vdots \\ x_N^{(k-1)} \end{bmatrix} + \begin{bmatrix} 0 & 0 & \cdots & \cdots & 0 \\ \times & 0 & 0 & \cdots & \vdots \\ a_{31} & a_{32} & 0 & 0 & \vdots \\ \vdots & \vdots & \times & 0 & 0 \\ \times & \cdots & \cdots & \times & 0 \end{bmatrix} \begin{bmatrix} x_1^{(k)} \\ \vdots \\ x_3^{(k)} \\ x_4^{(k-1)} \\ \vdots \\ x_N^{(k-1)} \end{bmatrix} \\ + \begin{bmatrix} 0 & \times & \cdots & \cdots & \cdots & \times \\ 0 & 0 & \times & \cdots & \cdots & \cdots \\ \vdots & 0 & 0 & a_{34} & a_{35} & \cdots \\ \vdots & \cdots & 0 & 0 & \times & \cdots \\ \vdots & \cdots & \cdots & 0 & 0 & \times \\ 0 & \cdots & \cdots & \cdots & 0 & 0 \end{bmatrix} \begin{bmatrix} x_1^{(k)} \\ \vdots \\ x_3^{(k)} \\ x_4^{(k-1)} \\ \vdots \\ x_N^{(k-1)} \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_3 \\ \vdots \\ \vdots \\ b_N \end{bmatrix} \quad (4.7)$$

Now using the decomposition $A = D + L + U$, this equation can be viewed as a particular case ($i = 3$) of the following general form

$$D \begin{bmatrix} x_1^{(k)} \\ \vdots \\ x_3^{(k)} \\ x_4^{(k)} \\ \vdots \\ x_N^{(k)} \end{bmatrix} + L \begin{bmatrix} x_1^{(k)} \\ \vdots \\ x_3^{(k)} \\ x_4^{(k)} \\ \vdots \\ x_N^{(k)} \end{bmatrix} + U \begin{bmatrix} x_1^{(k-1)} \\ \vdots \\ x_3^{(k-1)} \\ x_4^{(k-1)} \\ \vdots \\ x_N^{(k-1)} \end{bmatrix} = \begin{bmatrix} b_1 \\ \vdots \\ b_3 \\ \vdots \\ \vdots \\ b_N \end{bmatrix} \quad (4.8)$$

or in more compact form

$$D \vec{x}^{(k)} + L \vec{x}^{(k)} + U \vec{x}^{(k-1)} = \vec{b} \quad (4.9)$$

which is equivalent to

$$\vec{x}^{(k)} = (D + L)^{-1}(\vec{b} - U\vec{x}^{(k-1)}) \quad (4.10)$$

Recall that any iterative method for (4.6) can be expressed by

$$\vec{x}^{(k)} = B\vec{x}^{(k-1)} + \vec{c} \quad (4.11)$$

In summary, the stationary iterative methods so far are

$$\begin{array}{ll} \text{JM:} & B = -D^{-1}(L + U), \quad \vec{c} = D^{-1}\vec{b} \\ \text{GS:} & B = -(D + L)^{-1}U, \quad \vec{c} = (D + L)^{-1}\vec{b} \end{array} \quad (4.12)$$

Algorithm GS: Gauss-Seidel Method Solve $A\vec{x} = \vec{b}$.

Input: N : Number of unknowns and equations; a_{ij} : Entries of A , $i, j = 1 \cdots N$; b_i : Entries of \vec{b} , $i = 1 \cdots N$; TOL: Error Tolerance.

Output: $x_i^{(k)}$: Entries of $\vec{x}^{(k)}$ (approximate solution) or Error Message.

Step 1. Choose an initial guess $\vec{x}^{(0)}$ to the solution \vec{x} .

Step 2. For $k = 1, 2, 3 \cdots, k_{\max}$

Step 3. For $i = 1, 2, \cdots, N$

Step 4. $\sigma = 0$

Step 5. For $j = 1, 2, \cdots, i - 1$

Step 6. $\sigma = \sigma + a_{ij}x_j^{(k)}$

Step 7. End j loop

Step 8. For $j = i + 1, \dots, N$

Step 9. $\sigma = \sigma + a_{ij}x_j^{(k-1)}$

Step 10. End j loop

Step 11. $x_i^{(k)} = (b_i - \sigma)/a_{ii}$

Step 12. End i loop

Step 13. If $\|\vec{r}^{(k)}\|_\infty < \text{TOL} = 10^{-6}$ then Stop otherwise Set $\vec{x}^{(k-1)} = \vec{x}^{(k)}$ and Go To Step 2.

Step 14. End k loop

Step 15. Error: Not convergent with the max number of iterations k_{\max} and TOL.

$\vec{x}^{(k)} \approx \vec{x}$: An approximate solution.

$\vec{r}^{(k)} = A\vec{x}^{(k)} - \vec{b}$: Residual vector.

$E^{\vec{x}} = \|\vec{r}^{(k)}\|_\infty := \max_{1 \leq i \leq N} |r_i^{(k)}|$: Residual error in maximum norm.

Project 4.1. Consider the 1D Poisson Problem (1.1) (with $f(x) = 2$, $g_D = 0$, and $g_N = 0$) and implement the methods FDM and GS.

Input: N , A , \vec{b} , k_{\max} , TOL (write the input in the program).

Output:

N	k	$E^{\vec{x}}$	E^u	α
5				
9				
17				
33				
65				
129				